

# Proof Hints for Event-B

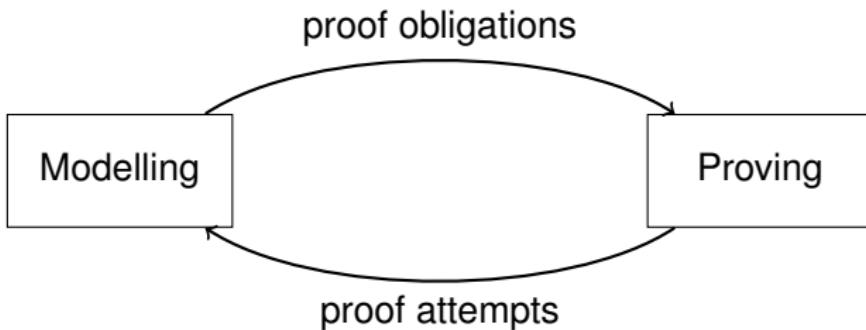
Thai Son Hoang

Institute of Information Security  
Swiss Federal Institute of Technology Zürich (ETH Zürich)

Rodin Workshop, Fontainebleau, France  
28th-29th February 2012



# Developing in Event-B. Modelling and Proving



- Proof obligations are generated from formal models.
- Failed proof attempts required to the models to be fixed.
- How about successful attempts, in particular, interactive proofs?

- Maintenance of interactive proofs is difficult.
- Better rate of automatic proofs
  - Better automatic provers (Isabelle, SMT)
  - Better proof profiles.
  - This talk: “Improve” the existing model.

## Idea

Expose more proof information in the model: “proof hints”



# Existing Proof Hints in Event-B/Rodin Platform

- **Theorems** (add hypothesis)
- **Witnesses** (existential instantiation)
- **Guard selection** (select hypotheses)



# Hypotheses Selection (1/2)

**invariants:**

**inv1** :  $x \in \mathbb{N}$

**inv2** :  $x \neq 0 \Rightarrow y \in \mathbb{N}$

**set**

**when**

**grd1** :  $x \in \{1, 2\}$

**then**

**act1** :  $x := y + 1$

**end**

**inv2**

**inv1**

**grd1**

$\vdash$

Modified **inv1**

$x \neq 0 \Rightarrow y \in \mathbb{N}$

$x \in \mathbb{N}$

$x \in \{1, 2\}$

$\vdash$

$y + 1 \in \mathbb{N}$

set/**inv1**/INV

- Selected hypotheses: **inv1** and **grd1**
- inv2** is required, added as a guard theorem.



# Hypotheses Selection (1/2)

invariants:

inv1 :  $x \in \mathbb{N}$

inv2 :  $x \neq 0 \Rightarrow y \in \mathbb{N}$

set

when

grd1 :  $x \in \{1, 2\}$

then

act1 :  $x := y + 1$

end

inv2

inv1

grd1

⊤

Modified inv1

$x \neq 0 \Rightarrow y \in \mathbb{N}$

$x \in \mathbb{N}$

$x \in \{1, 2\}$

⊤

$y + 1 \in \mathbb{N}$

set/inv1/INV

- Selected hypotheses: **inv1** and **grd1**
- inv2** is required, added as a guard theorem.



# Hypotheses Selection (1/2)

**invariants:**

**inv1** :  $x \in \mathbb{N}$

**inv2** :  $x \neq 0 \Rightarrow y \in \mathbb{N}$

**set**

**when**

**grd1** :  $x \in \{1, 2\}$

**then**

**act1** :  $x := y + 1$

**end**

**inv2**

**inv1**

**grd1**

$\vdash$

Modified **inv1**

$x \neq 0 \Rightarrow y \in \mathbb{N}$

$x \in \mathbb{N}$

$x \in \{1, 2\}$

$\vdash$

$y + 1 \in \mathbb{N}$

set/**inv1**/INV

- Selected hypotheses: **inv1** and **grd1**
- inv2** is required, added as a guard theorem.



# Hypotheses Selection (1/2)

invariants:

inv1 :  $x \in \mathbb{N}$

inv2 :  $x \neq 0 \Rightarrow y \in \mathbb{N}$

set

when

grd1 :  $x \in \{1, 2\}$

thm1 :  $x \neq 0 \Rightarrow y \in \mathbb{N}$

then

act1 :  $x := y + 1$

end

inv2

inv1

grd1

⊢

Modified inv1

$x \neq 0 \Rightarrow y \in \mathbb{N}$

$x \in \mathbb{N}$

$x \in \{1, 2\}$

⊢

$y + 1 \in \mathbb{N}$

set/inv1/INV

- Selected hypotheses: **inv1** and **grd1**
- inv2** is required, added as a guard theorem.



# Hypotheses Selection (2/2)

```
set
```

```
when
```

```
  grd1 :  $x \in \{1, 2\}$ 
```

```
  thm1 :  $x \neq 0 \Rightarrow y \in \mathbb{N}$ 
```

```
then
```

```
  act1 :  $x := y + 1$ 
```

```
end
```

```
set
```

```
when
```

```
  grd1 :  $x \in \{1, 2\}$ 
```

```
then
```

```
  act1 :  $x := y + 1$ 
```

```
select
```

```
  inv2
```

```
end
```

Cons for using theorem

- Copy/paste.
- An extra proof obligation (trivially discharged).



# Hypotheses Selection (2/2)

```
set
```

```
when
```

```
  grd1 :  $x \in \{1, 2\}$ 
```

```
  thm1 :  $x \neq 0 \Rightarrow y \in \mathbb{N}$ 
```

```
then
```

```
  act1 :  $x := y + 1$ 
```

```
end
```

```
set
```

```
when
```

```
  grd1 :  $x \in \{1, 2\}$ 
```

```
then
```

```
  act1 :  $x := y + 1$ 
```

```
select
```

```
inv2
```

```
end
```

Cons for using theorem

- Copy/paste.
- An extra proof obligation (trivially discharged).



# Do Case (1/3)

**invariants:**

**inv1** :  $a \leq c$

**inv2** :  $a \neq 1 \Rightarrow b = a + 1$

**inv3** :  $a = 1 \Rightarrow b \leq c$

set

**begin**

$a := b - 1$

**end**

$a \leq c$

$a \neq 1 \Rightarrow b = a + 1$

$a = 1 \Rightarrow b \leq c$

$\vdash$

$b - 1 \leq c$

set/**inv1**/INV

- Proof by cases:

- $a = 1$
- $a \neq 1$



# Do Case (2/3)

```
seta
when
  a = 1
then
  a := b - 1
end
```

```
setb
when
  a ≠ 1
then
  a := b - 1
end
```

```
set
refines  seta, setb
begin
  a := b - 1
end
```

- Duplication of proof obligations.
- Artificial merging step.



# Do Case (3/3)

```
set
begin
  a := b
  case-split
    a = 1 for inv1
end
```



# Summary

set

**when**

**grd1** :  $x \in \{1, 2\}$

**then**

**act1** :  $x := y + 1$

**select**

**inv2**

**end**

set

**begin**

$a := b$

**case-split**

$a = 1$  for **inv1**

**end**

- Using information of interactive proofs to “improve” the model.
- **Hints** (proof information) help with **proof automation**.
- Hints help to **understand model better**.
- How far should we go  
    in terms of exposing proof information in the model?
- A plug-in (a reasoner) that “interprete” proof hints.

